

# About convex optimization with constraints

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# Reminder: optimisation with constraints

## General case

$$\left\{ \begin{array}{ll} \text{Minimize} & f_0(\mathbf{z}) \quad \mathbf{z} \in \text{domain } \mathcal{D} \\ \text{s.t.} & f_i(\mathbf{z}) \leq 0 \quad i = 1 \dots m \\ & h_j(\mathbf{z}) = 0 \quad j = 1 \dots p \end{array} \right. \quad (1)$$

Notation:  $\mathbf{z}^*$  the solution of Pb (??)  
 $p^* = f_0(\mathbf{z}^*)$

## Define the Lagrangian as:

$$L(\mathbf{z}, \lambda, \nu) = f_0(\mathbf{z}) + \sum_{i=1}^m \lambda_i f_i(\mathbf{z}) + \sum_{j=1}^p \nu_j h_j(\mathbf{z})$$

with  $\lambda_i \geq 0$

## Reminder, 2

Define the Lagrange dual function as:

$$g(\lambda, \nu) = \inf\{L(\mathbf{z}, \lambda, \nu), \mathbf{z} \in \mathcal{D}, \lambda \geq 0\} \quad (2)$$

Notation:  $\lambda^*, \nu^* = \operatorname{argmin} g(\lambda, \nu)$   
 $d^* = g(\lambda^*, \nu^*)$

### Intuition

- ▶  $\lambda_i, \nu_j$  are the penalties paid for violating constraints
- ▶ For  $\mathbf{z}$  feasible (i.e.  $f_i(\mathbf{z}) \leq 0$  and  $h_j(\mathbf{z}) = 0$ ),

$$L(\mathbf{z}, \lambda, \nu) \leq f_0(\mathbf{z})$$

- ▶ Hence  $g(\lambda, \nu)$  provides a lower bound on the optimum  $p^*$ .
- ▶  $L$  is linear in  $\lambda$  and  $\nu$ : its minimum is very easy to compute;
- ▶  $g$  is concave in  $\lambda$  and  $\nu$  even if  $f_0, f_i$  and  $h_j$  are not convex in  $\mathbf{z}$ .

## Reminder, 3

With

$$g(\lambda, \nu) = \inf \left\{ f_0(\mathbf{z}) + \sum_{i=1}^m \lambda_i f_i(\mathbf{z}) + \sum_{j=1}^p \nu_j h_j(\mathbf{z}), \mathbf{z} \in \mathcal{D}, \lambda \geq 0 \right\}$$

**Then**

- ▶ Weak duality (always true)

$$d^* \leq p^*$$

- ▶ Strong duality: when  $d^* = p^*$ .  
**usually** holds for convex functions  $f_0$  and  $f_i$ .

## Reminder, 4

### Constraint qualifications

= conditions guaranteeing strong duality in convex pbs.

### Slater's conditions

if there exists  $\mathbf{z}$  s.t.  $f_i(\mathbf{z}) < 0$ , then strong duality

**Remark:** if strong duality then

$$\begin{aligned} f_0(\mathbf{z}^*) &= \inf_{\mathbf{z}} \left\{ f_0(\mathbf{z}) + \sum_{i=1}^m \lambda_i^* f_i(\mathbf{z}) + \sum_{j=1}^p \nu_j^* h_j(\mathbf{z}) \right\} \\ &\leq f_0(\mathbf{z}^*) + \sum_{i=1}^m \lambda_i^* f_i(\mathbf{z}^*) + \sum_{j=1}^p \nu_j^* h_j(\mathbf{z}^*) \\ &\leq f_0(\mathbf{z}^*) \end{aligned}$$

Therefore for each  $i = 1 \dots m$ ,  $\lambda_i^* f_i(\mathbf{z}^*) = 0$ :

if  $\lambda_i^* > 0$ ,  $f_i(\mathbf{z}^*) = 0$  (saturated or tight constraint).

if  $f_i(\mathbf{z}^*) < 0$  (loose constraint),  $\lambda_i^* = 0$ .

## Reminder, 5

**Karun-Kush-Tucker conditions** on  $\tilde{\mathbf{z}}, \tilde{\lambda}, \tilde{\nu}$

- ▶ Stationarity condition:  $f_i, h_j$  differentiable

$$\frac{\partial L(\mathbf{z}, \lambda, \nu)}{\partial \mathbf{z}} \Big|_{\mathbf{z}=\tilde{\mathbf{z}}} = 0$$

- ▶ Slackness condition:

$$\tilde{\lambda}_i f_i(\tilde{\mathbf{z}}) = 0 \text{ for } i = 1 \dots m$$

- ▶ Feasibility conditions

$$f_i(\tilde{\mathbf{z}}) \leq 0; \quad \tilde{\lambda}_i \geq 0$$

### Theorem

$\Rightarrow$  If  $\mathbf{z}^*$  is optimum of Pb (??), and  $\lambda^*, \nu^*$  minimize  $g$ , then KKT conditions are satisfied for  $(\mathbf{z}^*, \lambda^*, \nu^*)$ .

$\Leftarrow$  If  $f_i$  convex and  $h_j$  affine, then if KKT conditions are satisfied for  $\tilde{\mathbf{z}}, \tilde{\lambda}, \tilde{\nu}$ , then  $\tilde{\mathbf{z}}$  is optimum for Pb (??) and  $\tilde{\lambda}, \tilde{\nu}$  minimize  $g$ .